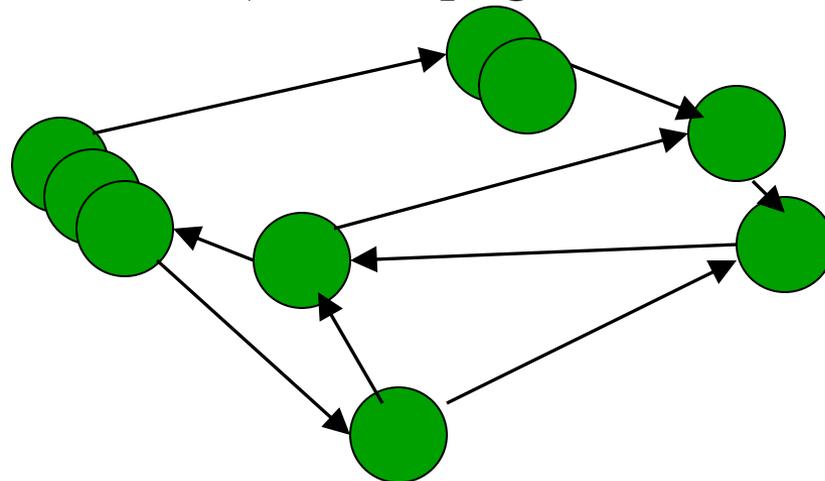


LINK ANALYSIS

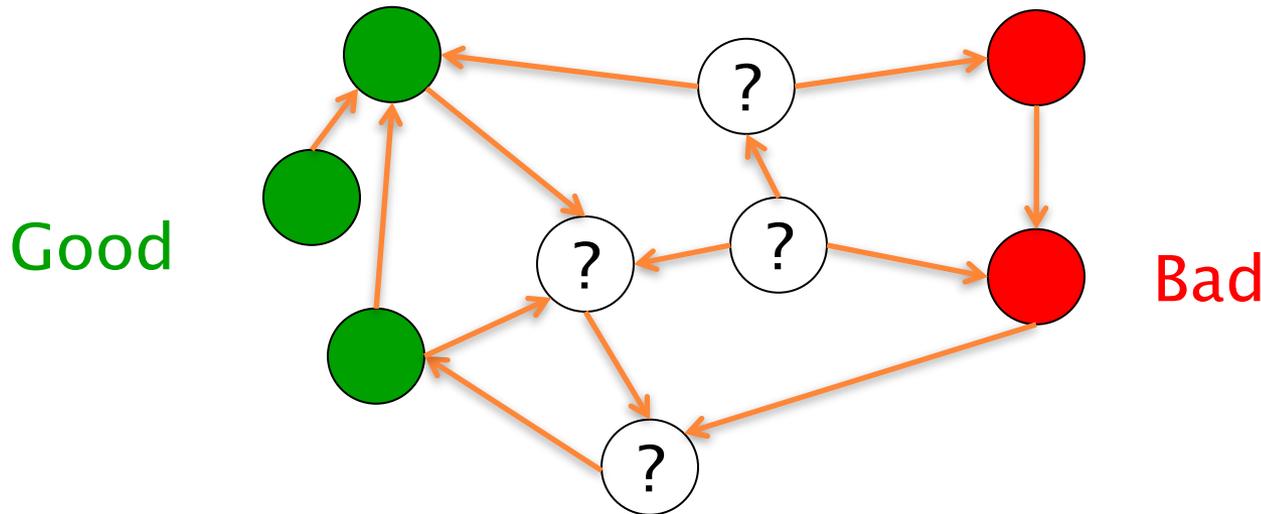
TODAY'S LECTURE – HYPERTEXT AND LINKS

- We look beyond the *content* of documents
 - We begin to look at the hyperlinks between them
- Address questions like
 - Do the links represent a conferral of authority to some pages? Is this useful for ranking?
 - How likely is it that a page pointed to by the CCF (Chinese Computer Federation) home page is about IT
- Big application areas
 - The Web
 - Email
 - Social networks



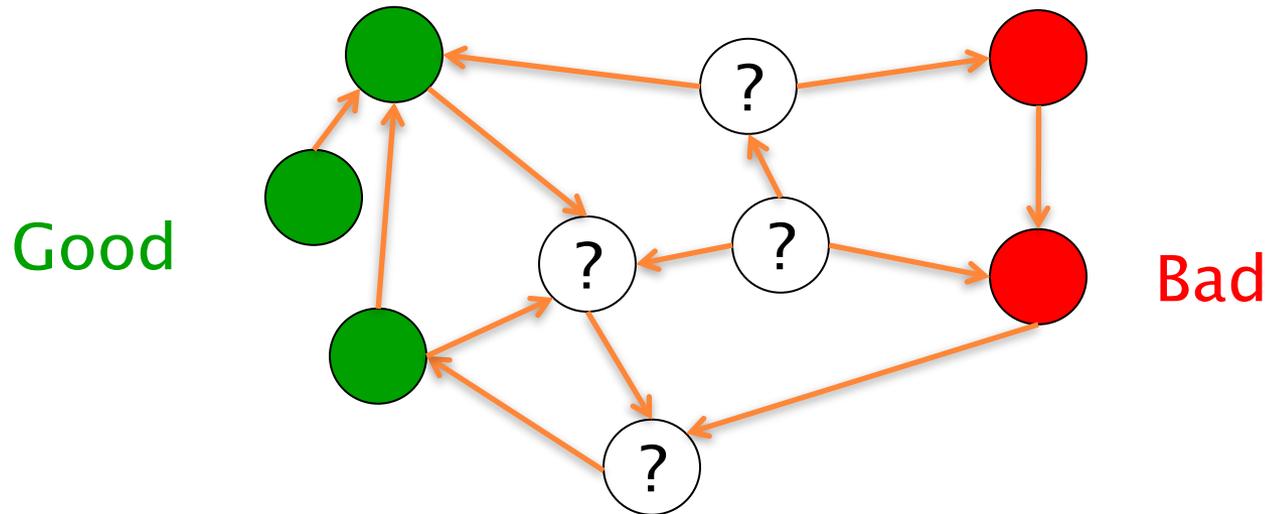
LINKS ARE EVERYWHERE

- Powerful sources of authenticity and authority
 - Mail spam – which email accounts are spammers?
 - Host quality – which hosts are “bad”?
 - Phone call logs
- The **Good**, The **Bad** and The Unknown



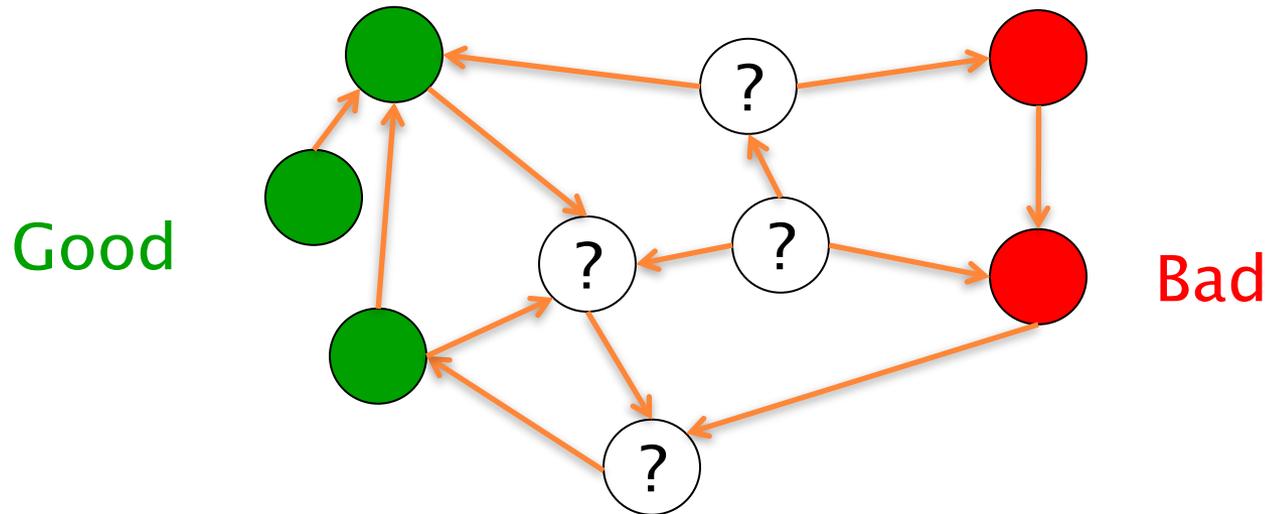
EXAMPLE 1: GOOD/BAD/UNKNOWN

- The Good, The Bad and The Unknown
 - Good nodes won't point to Bad nodes
 - All other combinations plausible



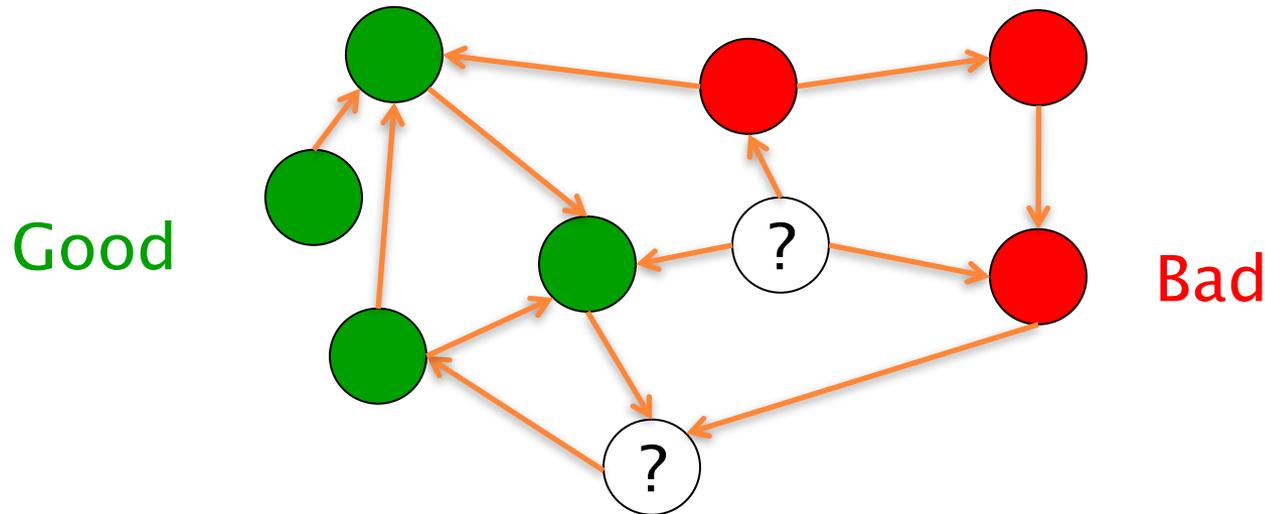
SIMPLE ITERATIVE LOGIC

- **Good** nodes won't point to **Bad** nodes
 - If you point to a **Bad** node, you're **Bad**
 - If a **Good** node points to you, you're **Good**



SIMPLE ITERATIVE LOGIC

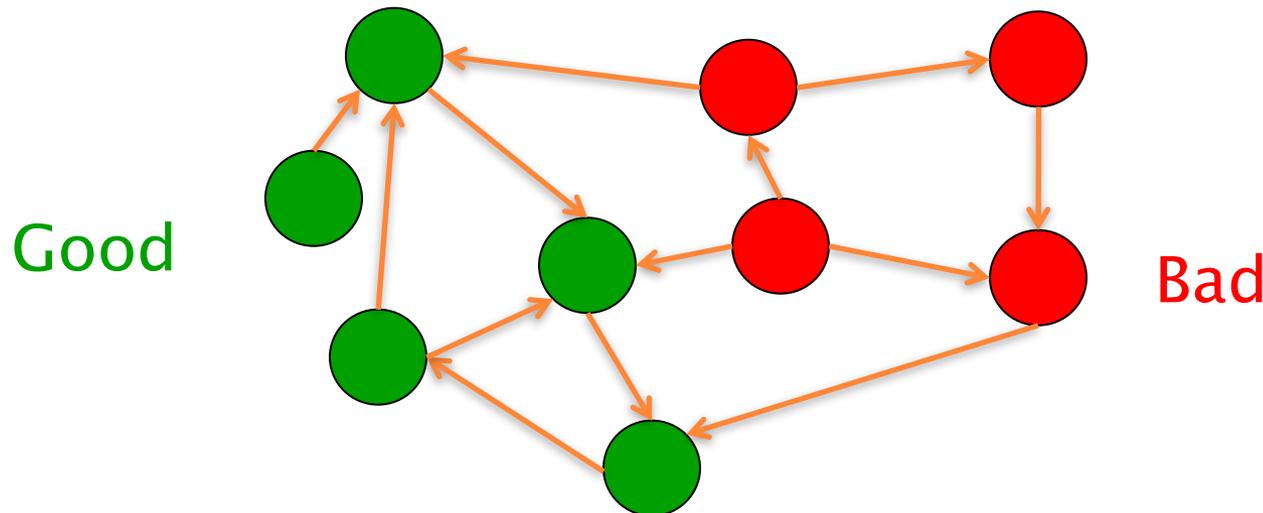
- **Good** nodes won't point to **Bad** nodes
 - If you point to a **Bad** node, you're **Bad**
 - If a **Good** node points to you, you're **Good**



SIMPLE ITERATIVE LOGIC

- **Good** nodes won't point to **Bad** nodes
 - If you point to a **Bad** node, you're **Bad**
 - If a **Good** node points to you, you're **Good**

Sometimes need **probabilistic** analogs – e.g., mail spam

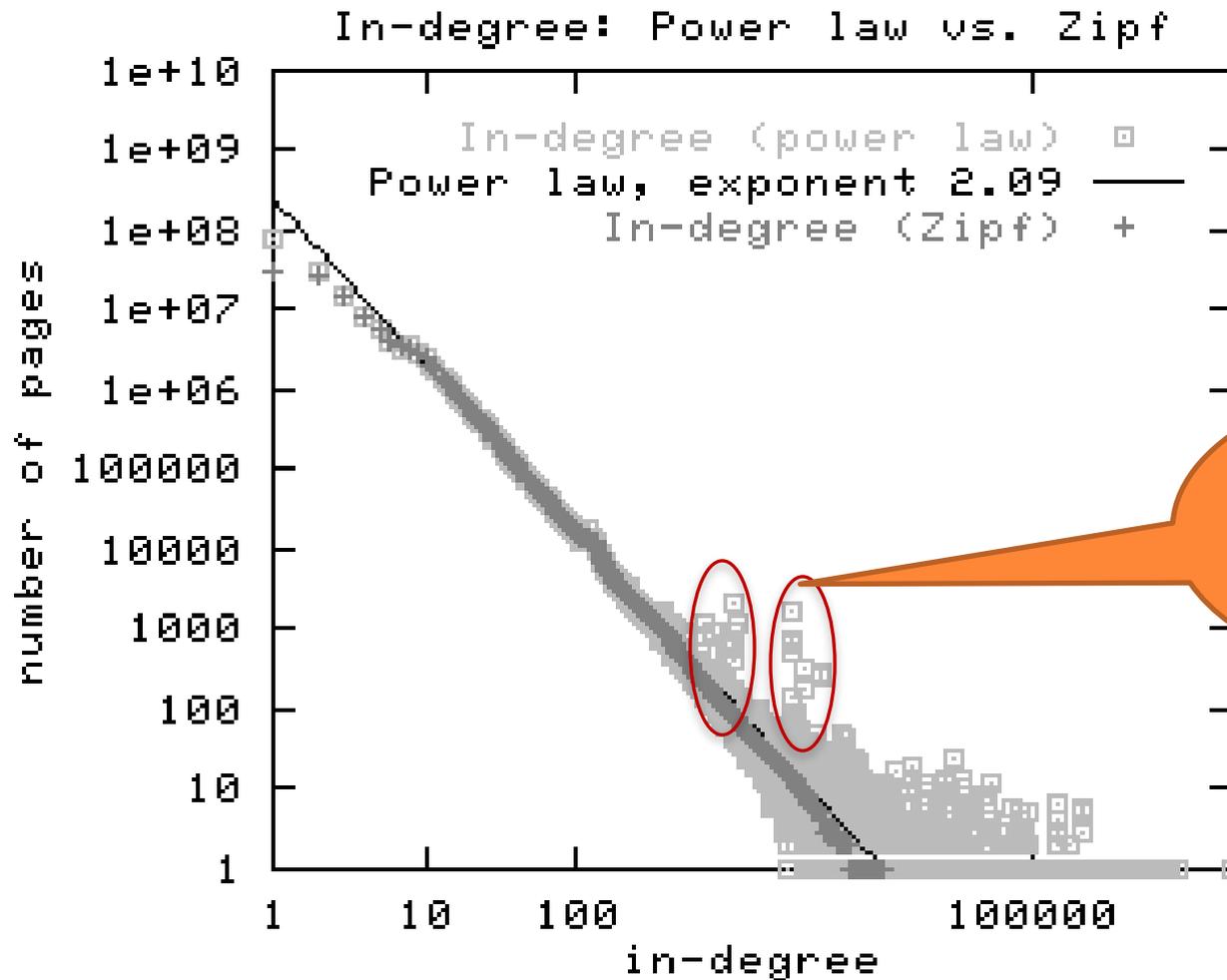


QUIZ: “GOOD NODES WON’T POINT TO BAD NODES?”

Can you think of one scenario in which this assumption doesn’t work?

EXAMPLE 2:

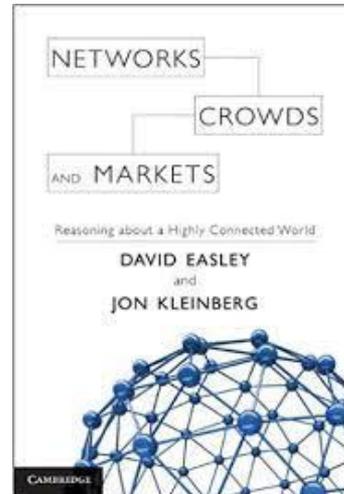
IN-LINKS TO PAGES – UNUSUAL PATTERNS 😊



Spammers
violating power
laws!

MANY OTHER EXAMPLES OF LINK ANALYSIS

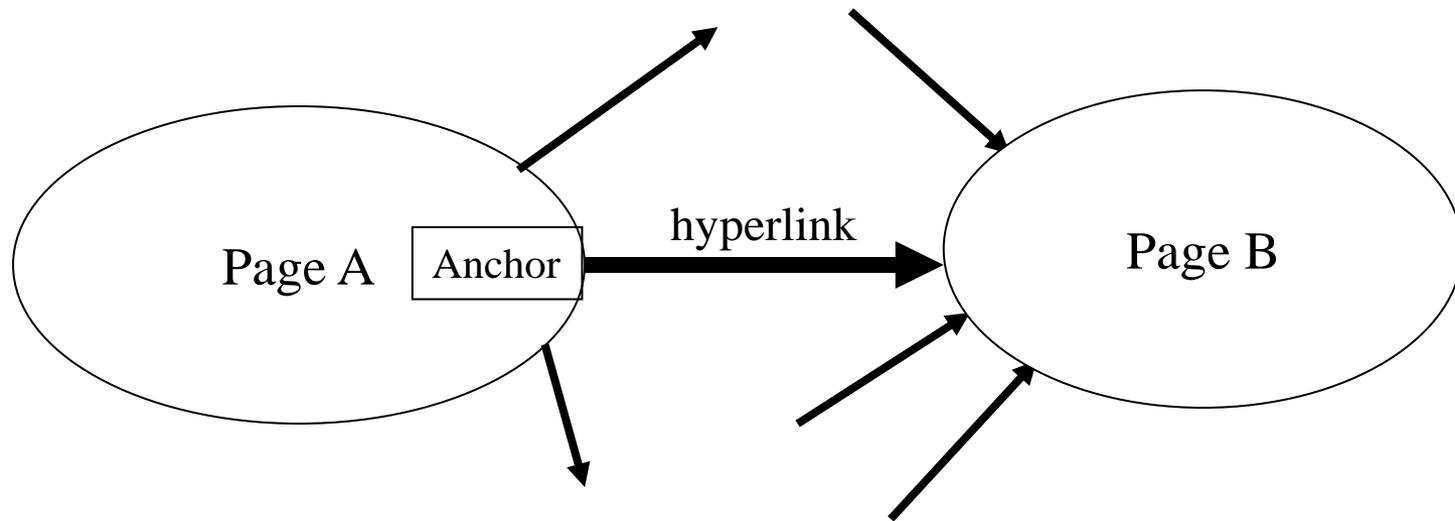
- Social networks are a rich source of grouping behavior
- E.g., Shoppers' affinity – Goel+Goldstein 2010
 - Consumers whose friends spend a lot, spend a lot themselves
- <http://www.cs.cornell.edu/home/kleinber/networks-book/> (Nice book!)



OUR PRIMARY INTEREST IN THIS COURSE

- IR features thus far based purely on text
 - TF-IDF scores
 - Cosine similarity
- Link structure plays key role in
 - Authority/popularity of document
 - Link-based clustering – topical structure from links
 - Links as features in classification – documents that link to one another are likely to be on the same subject
 - Crawling – Based on the links seen, where do we crawl next?

THE WEB AS A DIRECTED GRAPH

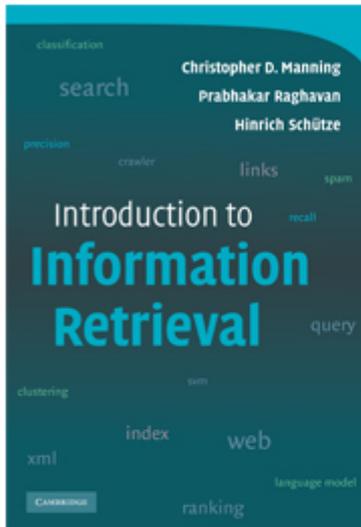


Assumption 1: A hyperlink between pages denotes a conferral of authority (quality signal)

Assumption 2: The text in the anchor of the hyperlink on page A describes the target page B

ASSUMPTION 1: REPUTED SITES

Introduction to Information Retrieval



This is the companion website for the following book.

[Christopher D. Manning](#), [Prabhakar Raghavan](#) and [Hinrich Schütze](#), *Introduction to Information Retrieval*

You can order this book at [CUP](#), at your local bookstore or on the internet. The best search

The book aims to provide a modern approach to information retrieval from a computer science perspective. It is available at [Cambridge University](#) and at the [University of Stuttgart](#).

We'd be pleased to get feedback about how this book works out as a textbook, what is missing, and what you think. Please send your comments to: [informationretrieval \(at\) yahoogroups \(dot\) com](mailto:informationretrieval@yahoogroups.com)

ASSUMPTION 2: ANNOTATION OF TARGET



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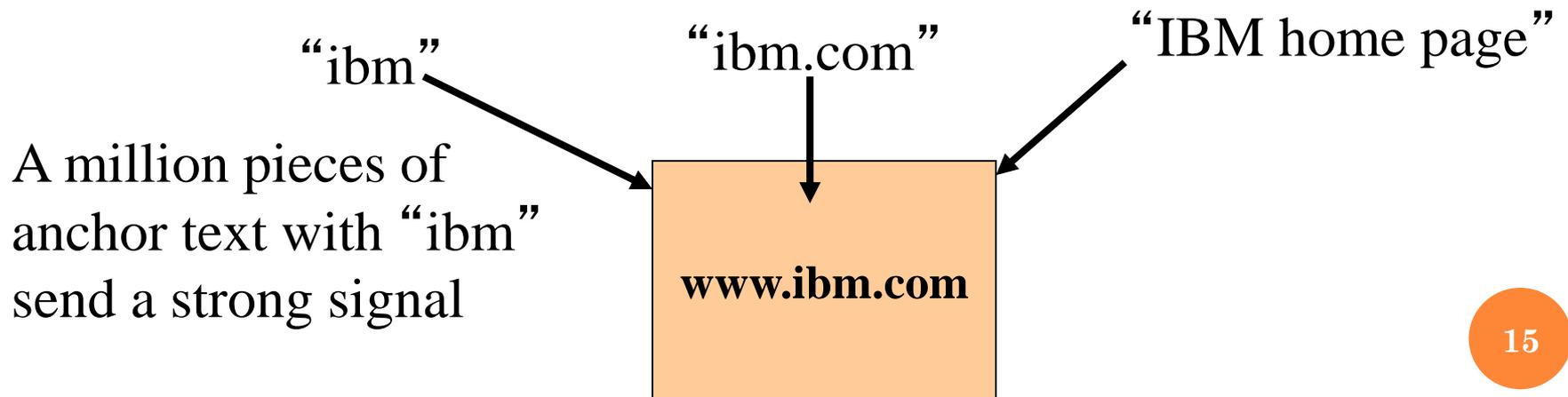
- Prospective Students
- General Public
- Corporations
- Alumni
- Current Students
- Faculty and Staff (Internal use)



ANCHOR TEXT

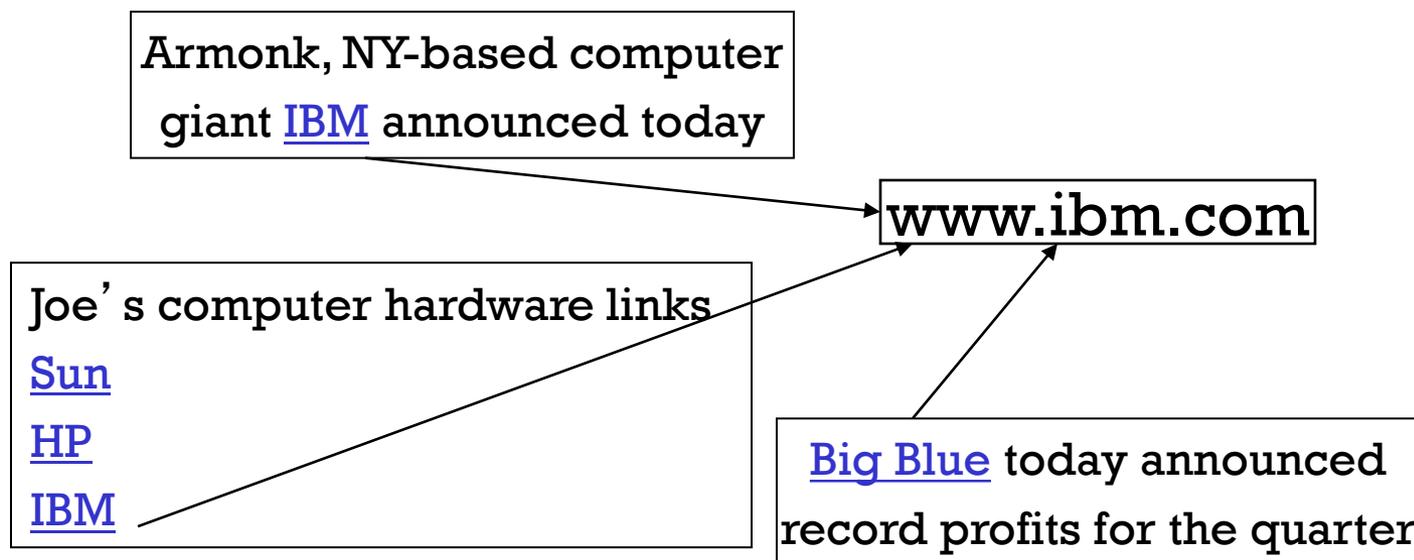
WWW WORM - MCBRYAN [MCBR94]

- For “*ibm*”, how to distinguish between:
 - IBM’s home page (mostly graphical)
 - IBM’s copyright page (high term freq. for ‘ibm’)
 - Rival’s spam page (arbitrarily high term freq.)



INDEXING ANCHOR TEXT

- When indexing a document D , include (with some weight) anchor text from links pointing to D .



INDEXING ANCHOR TEXT

- Can sometimes have unexpected effects, e.g., spam, **miserable failure**
- Can score anchor text with weight depending on the authority of the anchor page's website
 - E.g., if we were to assume that content from cnn.com or yahoo.com is authoritative, then trust (more) the anchor text from them
 - Increase the weight of off-site anchors (non-nepotistic scoring)

QUIZ

- If you want to defame Volkswagen (the German car maker) and claim that they only sell “lemons” (defective cars), with the help of anchor text, what would you do?

CONNECTIVITY SERVERS

CONNECTIVITY SERVER

- Support for fast queries on the web graph
 - Which URLs point to a given URL?
 - Which URLs does a given URL point to?
- Stores mappings in memory from
 - URL to outlinks (set of URLs), URL to inlinks (set of URLs)
- Applications
 - Link analysis
 - Web graph analysis
 - Connectivity, crawl optimization
 - Crawl control

BOLDI AND VIGNA 2004

- <http://www2004.org/proceedings/docs/1p595.pdf>
- Webgraph – set of algorithms and a java implementation
- Fundamental goal – maintain node adjacency lists in memory
 - For this, compressing the adjacency lists is the critical component

ADJACENCY LISTS

- The set of neighbors of a node
- Assume each URL represented by an integer
- E.g., for a 4 billion page web, need 32 bits per node
- Naively, this demands 64 bits to represent each hyperlink
- Boldi/Vigna get down to an average of ~ 3 bits/link
 - Further work achieves 2 bits/link

ADJACENCY LIST COMPRESSION

- Properties exploited in compression:
 - Similarity (between lists)
 - Locality (many links from a page go to “nearby” pages)
 - Use gap encodings in sorted lists
 - Distribution of gap values

MAIN IDEAS OF BOLDI/VIGNA

- Consider lexicographically ordered list of all URLs, e.g.,
 - www.stanford.edu/alchemy
 - www.stanford.edu/biology
 - www.stanford.edu/biology/plant
 - www.stanford.edu/biology/plant/copyright
 - www.stanford.edu/biology/plant/people
 - www.stanford.edu/chemistry

BOLDI/VIGNA

- Each of these URLs has an adjacency list
- Main idea: due to templates, the adjacency list of a node is similar to one of the 7 preceding URLs in the lexicographic ordering; 7 is a parameter
- Express adjacency list in terms of the “reference list”
- E.g., consider these adjacency lists

- 1, 2, 4, 8, 16, 32, 64

- 1, 4, 9, 16, 25, 36, 49, 64

- 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144

- 1, 4, 8, 16, 25, 36, 49, 64

Reference list

Pick a reference list that is the most similar to this list
 Encode as (-2), remove 9, add 8
 -2 means 2 lists before

QUIZ (BOLDI/VIGNA)

- Consider these adjacency lists
 - 1, 2, 4, 8, 16, 32, 64
 - 1, 4, 9, 16, 25, 36, 49, 64
 - 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144
 - 1, 4, 8, 16, 25, 36, 49, 64
 - The next list is encoded as
(-2), remove 1, add 35

What's the original adjacency list before encoding?

GAP ENCODINGS

- Given a sorted list of integers x, y, z, \dots , represent by $x, y-x, z-y, \dots$
- Compress each integer using a code
 - γ code - Number of bits = $1 + 2 \lfloor \lg x \rfloor$
 - δ code: ...
 - Information theoretic bound: $1 + \lfloor \lg x \rfloor$ bits
 - ζ code: Works well for integers from a power law
Boldi Vigna DCC 2004

MAIN ADVANTAGES OF BOLDI/VIGNA

- Depends only on locality in a canonical ordering
 - Lexicographic ordering works well for the web
- Adjacency queries can be answered very efficiently
 - To fetch out-neighbors, trace back **the chain of prototypes**
 - This chain is typically short in practice (since similarity is mostly intra-host)
 - Can also explicitly limit the length of the chain during encoding
- Easy to implement one-pass algorithm

LINK ANALYSIS: PAGERANK

CITATION ANALYSIS (INSPIRATION)

- Citation frequency
- **Bibliographic coupling frequency**
 - Articles that co-cite the same articles are related
- **Citation indexing**
 - Who is this author cited by? (Garfield 1972)
- Pagerank preview
 - Gabriel Pinski, Francis Narin:
Citation influence for journal aggregates of scientific publications: Theory, with application to the literature of physics. Inf. Process. Manag. **12(5)**: 297-312 (1976)
 - Asked: which journals are authoritative?

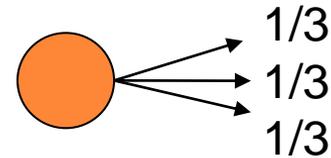
THE WEB ISN'T SCHOLARLY CITATION

- Millions of participants, each with self interests
- Spamming is widespread
- Once search engines began to use links for ranking (roughly 1998), link spam grew
 - You can join a *link farm* – a group of websites that heavily link to one another

PAGERANK SCORING

- Imagine a user doing a random walk on web pages:

- Start at a random page
- At each step, go out of the current page along one of the links on that page, equiprobably



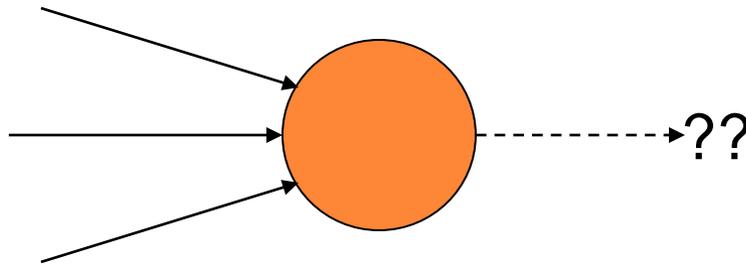
- “In the long run” each page has a long-term visit rate - use this as the page’s score.

QUIZ

- Can you think of a scenario where the probabilities of outgoing links are not equal?

NOT QUITE ENOUGH

- The web is full of dead-ends.
 - Random walk can get stuck in dead-ends.
 - Makes no sense to talk about long-term visit rates.



TELEPORTING

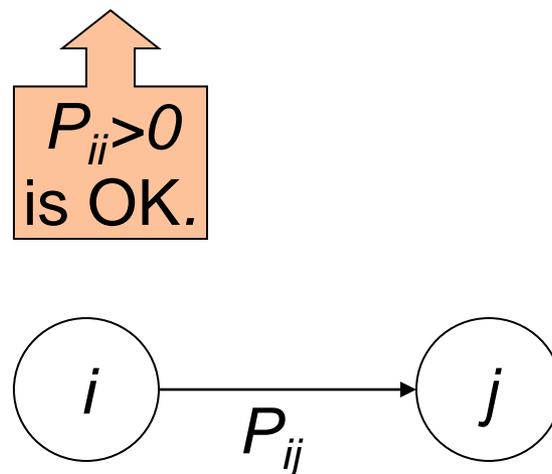
- At a dead end, jump to a random web page.
- At any non-dead end, with probability 10%, jump to a random web page.
 - With remaining probability (90%), go out on a random link.
 - 10% - a parameter.

RESULT OF TELEPORTING

- Now cannot get stuck locally.
- There is a long-term rate at which any page is visited (not obvious, will show this).
- How do we compute this visit rate?

MARKOV CHAINS

- A Markov chain consists of n states, plus an $n \times n$ transition probability matrix \mathbf{P} .
- At each step, we are in one of the states.
- For $1 \leq i, j \leq n$, the matrix entry P_{ij} tells us the probability of j being the next state, given we are currently in state i .

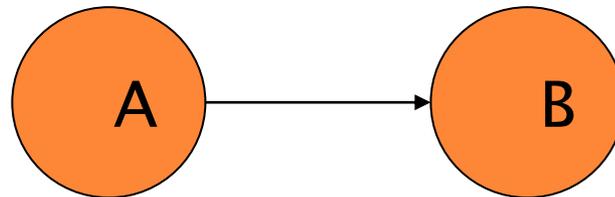


MARKOV CHAINS

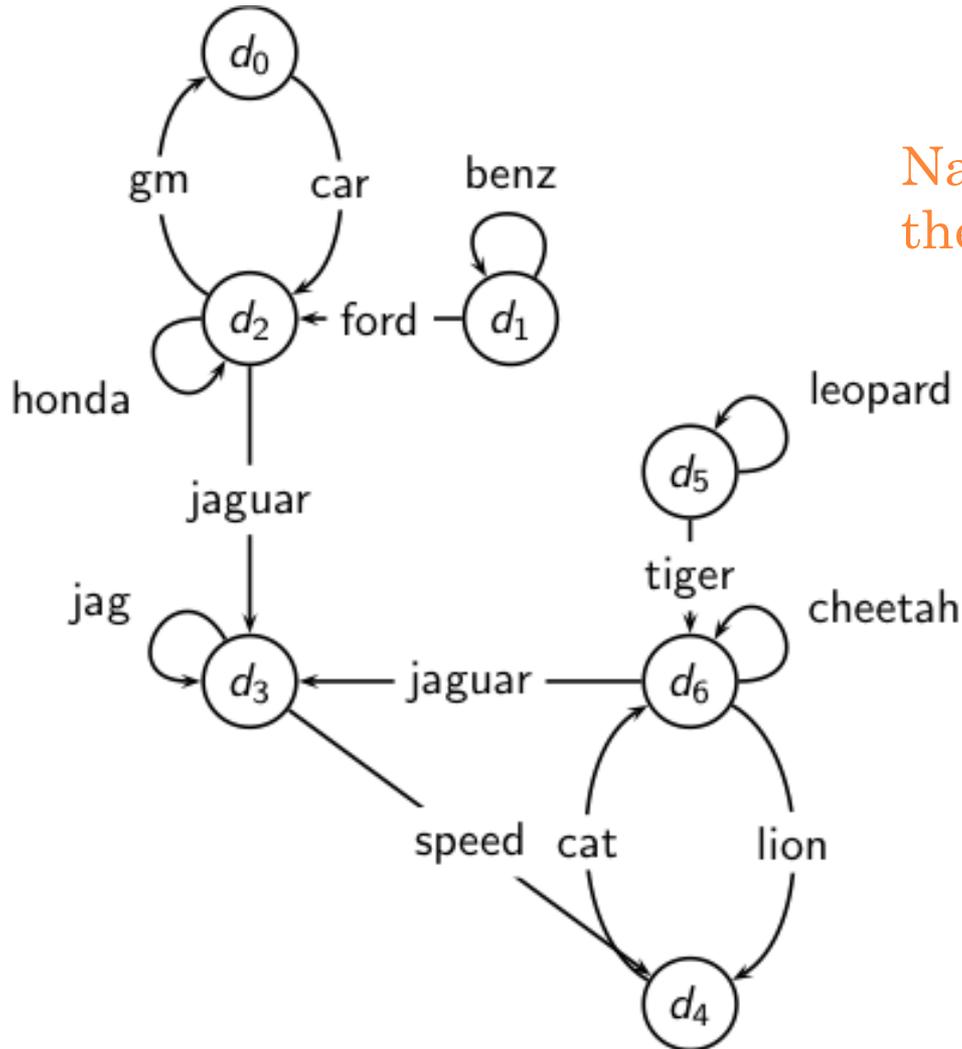
- Clearly, for all i , $\sum_{j=1}^n P_{ij} = 1$.
- Markov chains are abstractions of random walks.

QUIZ: MARKOV CHAIN

- The figure represents a random walk Markov chain with teleporting (100% probability for deadends and 10% probability for non-deadends). Beginning at state A, what is the probability of reaching state B after 2 steps?



EXAMPLE WEB GRAPH



Names on the edge are
the anchor texts

Link matrix for example

	d_0	d_1	d_2	d_3	d_4	d_5	d_6
d_0	0	0	1	0	0	0	0
d_1	0	1	1	0	0	0	0
d_2	1	0	1	1	0	0	0
d_3	0	0	0	1	1	0	0
d_4	0	0	0	0	0	0	1
d_5	0	0	0	0	0	1	1
d_6	0	0	0	1	1	0	1

Transition probability matrix P for example

	d_0	d_1	d_2	d_3	d_4	d_5	d_6
d_0	0.00	0.00	1.00	0.00	0.00	0.00	0.00
d_1	0.00	0.50	0.50	0.00	0.00	0.00	0.00
d_2	0.33	0.00	0.33	0.33	0.00	0.00	0.00
d_3	0.00	0.00	0.00	0.50	0.50	0.00	0.00
d_4	0.00	0.00	0.00	0.00	0.00	0.00	1.00
d_5	0.00	0.00	0.00	0.00	0.00	0.50	0.50
d_6	0.00	0.00	0.00	0.33	0.33	0.00	0.33

After normalization

ERGODIC MARKOV CHAINS

- For any *ergodic* Markov chain, there is a unique long-term visit rate for each state.
 - *Steady-state probability distribution.*
- Over a long time-period, we visit each state in proportion to this rate.
- It doesn't matter where we start.

PROBABILITY VECTORS

- A probability (row) vector $\mathbf{x} = (x_1, \dots, x_n)$ tells us where the walk is at any point.
- E.g., $(000\dots 1\dots 000)$ means we're in state i .

$$1 \leq i \leq n$$

More generally, the vector $\mathbf{x} = (x_1, \dots, x_n)$ means the walk is in state i with probability x_i .

$$\sum_{i=1}^n x_i = 1.$$

CHANGE IN PROBABILITY VECTOR

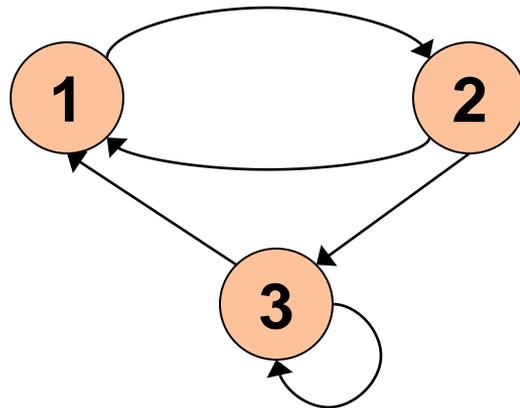
- If the probability vector is $\mathbf{x} = (x_1, \dots, x_n)$ at this step, what is it at the next step?
- Recall that row i of the transition prob. Transition matrix \mathbf{P} tells us where we go next from state i .
- So from \mathbf{x} , our next state is distributed as \mathbf{xP}
 - The one after that is \mathbf{xP}^2 , then \mathbf{xP}^3 , etc.
 - (Where) Does this converge?

HOW DO WE COMPUTE THIS VECTOR?

- Let $\mathbf{a} = (a_1, \dots, a_n)$ denote the row vector of steady-state probabilities.
- If our current position is described by \mathbf{a} , then the next step is distributed as \mathbf{aP} .
- But \mathbf{a} is the steady state, so $\mathbf{a}=\mathbf{aP}$.
- Solving this matrix equation gives us \mathbf{a} .
 - So \mathbf{a} is the (left) eigenvector for \mathbf{P} .
 - (Corresponds to the “principal” eigenvector of \mathbf{P} with the largest eigenvalue.)
 - Transition probability matrices always have largest eigenvalue 1.

QUIZ: PAGERANK

- Compute the steady-state probabilities of the following graph: (assuming equal probabilities for outgoing links of any node, and 0 transition probability if there's no link)



LINK ANALYSIS: HITS

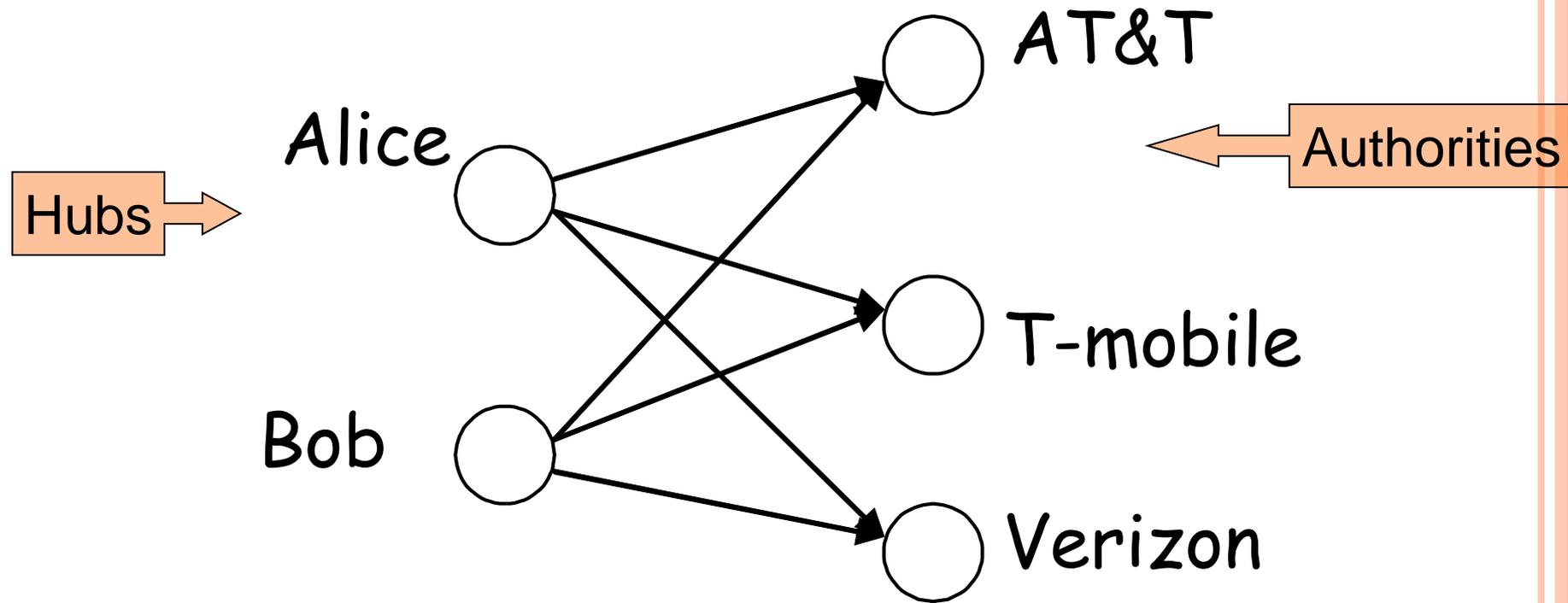
HYPERLINK-INDUCED TOPIC SEARCH (HITS)

- In response to a query, instead of an ordered list of pages each meeting the query, find two sets of inter-related pages:
 - *Hub pages* are good lists of links on a subject.
 - e.g., “Bob’s list of cancer-related links.”
 - *Authority pages* occur frequently on good hubs for the subject.
- Best suited for “broad topic” queries rather than for page-finding queries.
- Gets at a broader slice of common *opinion*.

HUBS AND AUTHORITIES

- Thus, a good hub page for a topic *points* to many authoritative pages for that topic.
- A good authority page for a topic is *pointed* to by many good hubs for that topic.
- Circular definition - will turn this into an iterative computation.

THE HOPE



Mobile telecom companies

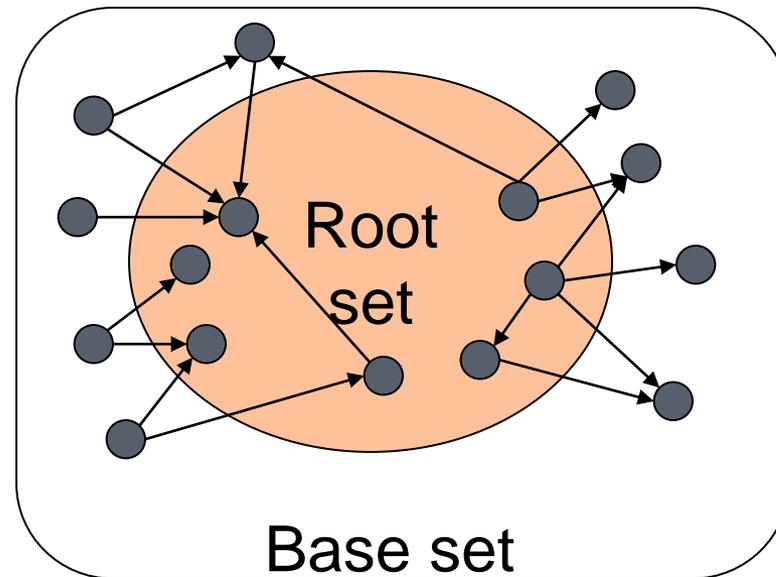
HIGH-LEVEL SCHEME

- Extract from the web a base set of pages that *could* be good hubs or authorities.
- From these, identify a small set of top hub and authority pages;
 - iterative algorithm.

BASE SET

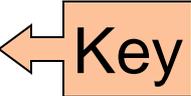
- Given text query (say *browser*), use a text index to get all pages containing *browser*.
 - Call this the root set of pages.
- Add in any page that either
 - points to a page in the root set, or
 - is pointed to by a page in the root set.
- Call this the base set.

VISUALIZATION



Get in-links (and out-links) from a *connectivity server*

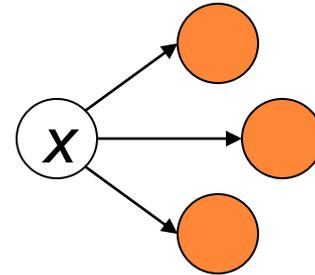
DISTILLING HUBS AND AUTHORITIES

- Compute, for each page x in the base set, a hub score $h(x)$ and an authority score $a(x)$.
- Initialize: for all x , $h(x) \leftarrow 1$; $a(x) \leftarrow 1$;
- Iteratively update all $h(x)$, $a(x)$;  Key
- After iterations
 - output pages with highest $h()$ scores as top hubs
 - highest $a()$ scores as top authorities.

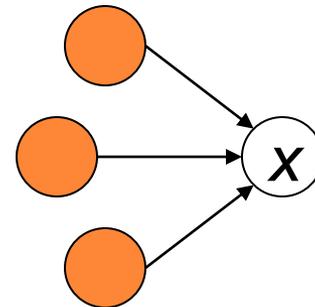
ITERATIVE UPDATE

- Repeat the following updates, for all x :

$$h(x) \leftarrow \sum_{x \mapsto y} a(y)$$



$$a(x) \leftarrow \sum_{y \mapsto x} h(y)$$



SCALING

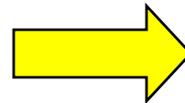
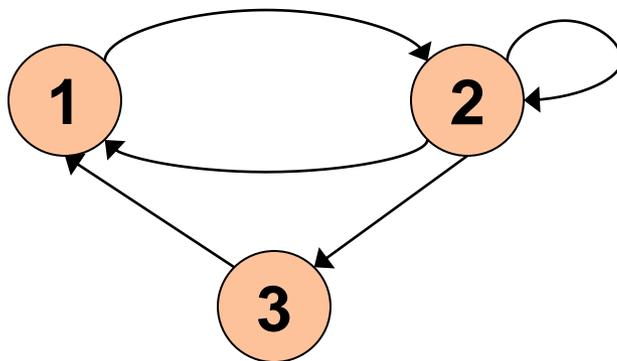
- To prevent the $h()$ and $a()$ values from getting too big, can scale down after each iteration.
- Scaling factor doesn't really matter:
 - we only care about the *relative* values of the scores.

HOW MANY ITERATIONS?

- Claim: relative values of scores will converge after a few iterations:
 - in fact, **suitably scaled**, $h()$ and $a()$ scores settle into a steady state!
 - L2-norm: normalize the $h()$ by the sqrt of sum of the squares of all $h()$ scores; and ditto for $a()$.
 - proof of this comes later.
- In practice, **~5 iterations get you close to stability.**

PROOF OF CONVERGENCE

- $n \times n$ adjacency matrix **A**:
 - each of the n pages in the base set has a row and column in the matrix.
 - Entry $A_{ij} = 1$ if page i links to page j , else $= 0$.
 - i is the row index and j is the col index



	1	2	3
1	0	1	0
2	1	1	1
3	1	0	0

HUB/AUTHORITY VECTORS

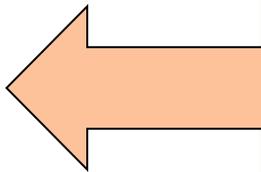
- View the hub scores $h()$ and the authority scores $a()$ as vectors with n components.
- Recall the iterative updates

$$h(x) \leftarrow \sum_{x \mapsto y} a(y)$$

$$a(x) \leftarrow \sum_{y \mapsto x} h(y)$$

REWRITE IN MATRIX FORM

- $\mathbf{h} = \mathbf{A}\mathbf{a}$.
- $\mathbf{a} = \mathbf{A}^t\mathbf{h}$.



Recall \mathbf{A}^t is the transpose of \mathbf{A} .
Change source to destination

Substituting, $\mathbf{h} = \mathbf{A}\mathbf{A}^t\mathbf{h}$ and $\mathbf{a} = \mathbf{A}^t\mathbf{A}\mathbf{a}$.

Thus, \mathbf{h} is an eigenvector of $\mathbf{A}\mathbf{A}^t$ and \mathbf{a} is an eigenvector of $\mathbf{A}^t\mathbf{A}$.

Further, our algorithm is a special case of well known algorithm for computing eigenvectors: the *power iteration* method (https://en.wikipedia.org/wiki/Power_iteration).



Guaranteed to converge.

ISSUES

- Topic Drift
 - Off-topic pages can cause off-topic “authorities” to be returned
 - E.g., the neighborhood graph can be about a “super topic” (larger, more general topic)
- Mutually Reinforcing Affiliates
 - Affiliated pages/sites can boost each others’ scores
 - Linkage between affiliated pages is not a useful signal

PageRank vs. HITS: Discussion

- PageRank can be precomputed, HITS has to be computed at query time.
- HITS is too expensive in most application scenarios.
- PageRank and HITS make two different design choices concerning (i) the eigenproblem formalization (ii) the set of pages to apply the formalization to.
- These two are orthogonal.
- We could also apply HITS to the entire web and PageRank to a small base set.
- Claim: On the web, a good hub almost always is also a good authority.
- The actual difference between PageRank ranking and HITS ranking is therefore not as large as one might expect.

RESOURCES

- IIR Chap 21
- <http://www2004.org/proceedings/docs/1p309.pdf>
- <http://www2004.org/proceedings/docs/1p595.pdf>
- <http://www2003.org/cdrom/papers/refereed/p270/kamvar-270-xhtml/index.html>
- <http://www2003.org/cdrom/papers/refereed/p641/xhtml/p641-mccurley.html>
- [The WebGraph framework I: Compression techniques \(Boldi et al. 2004\)](#)